



1. through 15. (Canceled)

16. (Currently Amended) A computer-implemented method for measuring a complexity of nested object state transition diagrams that are represented as data structures in the computer, the method comprising:

a) using [a] the computer, determining a plurality of graphs of object state transitions at K levels  $l_k$ , wherein:

(A)  $0 < k < K$ ;

(B) one or more of said graphs at level  $l_{k+1}$  are expansions of one or more of said graphs at level  $l_k$ ; and

(C) said graphs comprise a plurality of nodes to represent a corresponding plurality of states of use-cases and a plurality of edges to represent a corresponding plurality of transitions between the states; and

b) using the computer, determining a complexity for said plurality of graphs.

17. (Previously Presented) The computer-implemented method of claim 16, wherein:  
the state of an object at level  $k$  is a super state of an object state at level  $k+1$ , and  
the state of an object at level  $k$  is a sub state of an object state at level  $k-1$ , and  
the method further includes identifying the state transition relationship between super state and sub state as one of a group consisting of "in," "out," and "inout".

18. (Previously Presented) The computer-implemented method of claim 17, further comprising:

selecting said level  $l_k = l_K$ ;

expanding a graph  $L$  from said selected level  $l_k$  with at least one graph from said level  $l_{k+1}$ ;

determining paths in said expanded graph; and

determining a number of conditions in said transition paths.

19. (Previously Presented) The computer-implemented method of claim 18 further comprising:

removing an unnecessary path from said at least one graph.

20. (Previously Presented) The computer-implemented method of claim 19 wherein said one or more unnecessary paths comprise one or more numbers of a group consisting of:

(exception → fallout),

(exception → exception),

(one-repetition loop → fallout),

(one repetition loop → cancelled), and

(one repetition loop → exception).

21. (Previously Presented) ~~A computer-implemented method of claim 16, to~~ A computer-implemented method for measuring a complexity of nested object state transition diagrams that are represented as data structures in the computer, the complexity measuring method comprising:

a) using the computer, determining a plurality of graphs of object state transitions at K levels  $l_k$ , wherein:

(A)  $0 < k < K$ ;

(B) one or more of said graphs at level  $l_{k+1}$  are expansions of one or more of said graphs at level  $l_k$ ; and

(C) said graphs comprise a plurality of nodes to represent a corresponding plurality of states of use-cases and a plurality of edges to represent a corresponding plurality of transitions between the states; and

b) using the computer, determining a complexity for said plurality of graphs, the complexity determining step including measuring [e] a nested object state transition complexity between two super-states by recursively applying Equation 6:

$$STPC_{k,p,q} = \sum_{i=1}^{m_k} \left( \prod_{j=1}^{n_{k,i}} (C_{k,i,j} + (STPC_{k+1,i,j} - sub_{k,i,j})) + mul_{k,i}(N) \right)$$

Equation 6

wherein:

$m_k$  = a number of transition paths of level-k object ( $k>1$ ) between two super-states or a number of transition paths of level-1 object;

$n_{k,i}$  = a number of states along path  $i$  for level-k object;

$C_{k,i,j}$  is a number of conditions to bring level-k object state from  $Sk,i,j-1$  to  $Sk,i,j$  ;

$STPC_{k+1,i,j}$  = Substate Transition Complexity between state  $Sk,i,j-1$  and  $Sk,i,j$ ;

$sub_{k,i,j}=0$  if  $STPC_{k+1,i,j}=0$  and  $sub_{k,i,j}=1$  if ( $STPC_{k+1,i,j}\neq 0$ ) ;

if a multiplicity between level-k object and level- $(k+1)$  is  $1:N$ , then  $mul_{k,i}(N)=0$  if ( $N=1$ ) and  $mul_{k,i}(N)=r$  if ( $N>1$ ); and

$r$  is a number of times that level- $k+1$  objects transit back to level- $k$  in path  $i$ .

22. (Canceled)

*Add the following new claims 23-26:*

23. (New) The computer-implemented method of claim 21, wherein:

the state of an object at level  $k$  is a super state of an object state at level  $k+1$ , and

the state of an object at level  $k$  is a sub state of an object state at level  $k-1$ , and

the method further includes identifying the state transition relationship between super state and sub state as one of a group consisting of “in,” “out,” and “inout”.

24. (New) The computer-implemented method of claim 23, further comprising:

selecting said level  $l_k = l_k$ ;

expanding a graph  $L$  from said selected level  $l_k$  with at least one graph from said level  $l_{k+1}$ ;

determining paths in said expanded graph; and

determining a number of conditions in said transition paths.

25. (New) The computer-implemented method of claim 24 further comprising:

removing an unnecessary path from said at least one graph.

26. (New) The computer-implemented method of claim 25 wherein said one or more unnecessary paths comprise one or more numbers of a group consisting of:

(exception → fallout),

(exception → exception),

(one-repetition loop → fallout),

(one repetition loop → cancelled), and

(one repetition loop → exception).